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Quantum mutual information based on quantum conditional probability

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Plan of my talk

1. Motivation:
 - Simple review of the classical mutual entropy with three types of interpretation-

2. (1) Two types of quantum conditional probability
 - (2) Four scenarios of mutual entropy
 - (3) Their relations

3. The measurement-induced disturbance of the quantum channel Φ

4. Summary

Classical information theory

Composite system $x \otimes y = \{(x_i, y_j)\}$ with its probabilities

$$p^{xy} = (p_{ij}^{xy}), \quad p^x = (p_i^x), \quad p^y = (p_j^y), \quad p_{j|i}^y = (p_{j|i})$$

Bayes formula :

$$p_{j|i} = \frac{p_{ij}^{xy}}{p_i^x} \Leftrightarrow p_{ij}^{xy} = p_i^x p_{j|i}$$

$$S(x) = S(p^x) = -\sum_i p_i^x \log p_i^x$$

< Mutual entropy by relative entropy >

$$I(x, y) \equiv \sum p_{ij}^{xy} \log \frac{p_{ij}^{xy}}{p_i^x p_j^y} = S(x) + S(y) - S(x \otimes y) \quad (1)$$

$$= \sum p_i^x p_{j|i} \log \frac{p_{j|i}}{p_j^y} = S(y) - S(y|x) \quad (2)$$

Interpretation of mutual entropy:

(1) Correlation:

The common information encoded into x and y .

(2) Information gain by measurement:

$$S(y|x) \equiv \sum_i p_i^x S(p_{j|i}^y) = \sum_i p_i^x \left(-\sum_j p_{j|i} \log p_{j|i} \right)$$

The conditional entropy measures how uncertain we are of y on the average when we know x .

< Mutual entropy on communication scheme >

Input system: (x, p^x) , Output system: (y, p^y)

Channel: $T = (T_{ij}) = (p_{j|i})$ transition prob. matrix

$$p^y = T(p^x) \quad (\Leftrightarrow p_j^y = \sum_i p_i^x T_{ij})$$

$$I(x, y) = \sum_{ij} p_i^x p_{j|i} \log \frac{p_{j|i}}{\sum_k p_k^x p_{j|k}} = I(p^x; T) \quad (3)$$

$$= S(p^x) + S(T(p^x)) - S(p_T^{xy}) \quad (1)'$$

$$= S(T(p^x)) - \sum_i p_i^x S(p_{i|}^y) \quad (2)'$$

(3) Transmitted information through the channel:

The amount received less the part of this which is due to noise.

<Shannon's fundamental inequality>

$$0 \leq I(p^x; T) \leq \min \{ S(p^x), S(T(p^x)) \}$$

$I(x, y)$ has three types of interpretation at the same time based on Bayes formula

$$I(x, y) = S(x) + S(y) - S(x \otimes y)$$

$$= S(T(p^x)) - \sum_i p_i^x S(p_{i|}^y) = I(p^x; T)$$

How about quantum composite system?

Mutual entropies on a quantum composite system

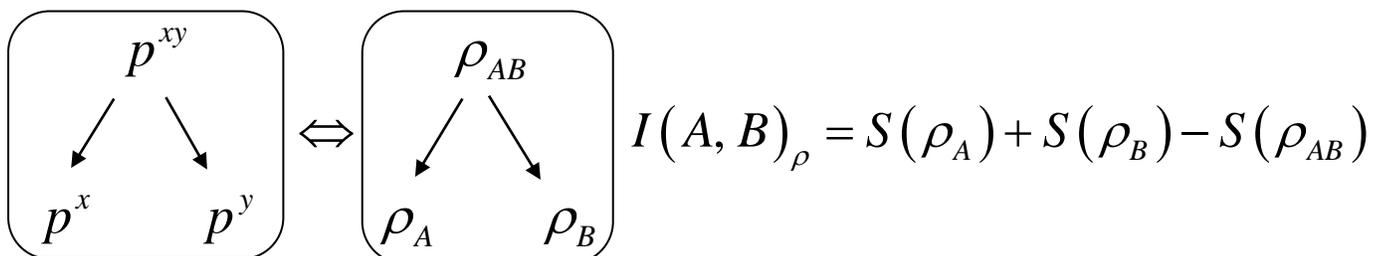
Composite system $H_A \otimes H_B$ with its density operators

$$\rho_{AB} \in S(H_{AB})$$

$$\text{Tr}_B \rho_{AB} = \rho_A \in S(H_A), \quad \text{Tr}_A \rho_{AB} = \rho_B \in S(H_B)$$

$$S(\rho_A) = -\text{tr} \rho_A \log \rho_A$$

<Quantum mutual entropy: scenario 1 > (by some authors)



A measure of a quantum correlation encoded into a compound state ρ_{AB}

(Via quantum relative entropy)

$$D(\rho \parallel \sigma) \equiv \text{Tr} \rho (\log \rho - \log \sigma)$$

(by Umegaki)

$$\begin{aligned} I(A, B)_\rho &\equiv D(\rho_{AB} \parallel \rho_A \otimes \rho_B) \\ &= \text{Tr}(\rho_{AB} \log \rho_{AB} - \log \rho_A \otimes \rho_B) \\ &= S(\rho_A) + S(\rho_B) - S(\rho_{AB}) \end{aligned}$$

$I(A, B)_\rho$ **has a same structure with** $I(x, y)$.

<Quantum mutual entropy via quantum channel:

scenario 2 >

The quantum channel Φ via Stinespring representation:

$$\Phi(\rho_A) = \text{Tr}_E V \rho_A V^*$$

where $V: H_A \rightarrow H_B \otimes E$ is a linear isometry

and E may be interpreted as an environment Hilbert space.

⇓

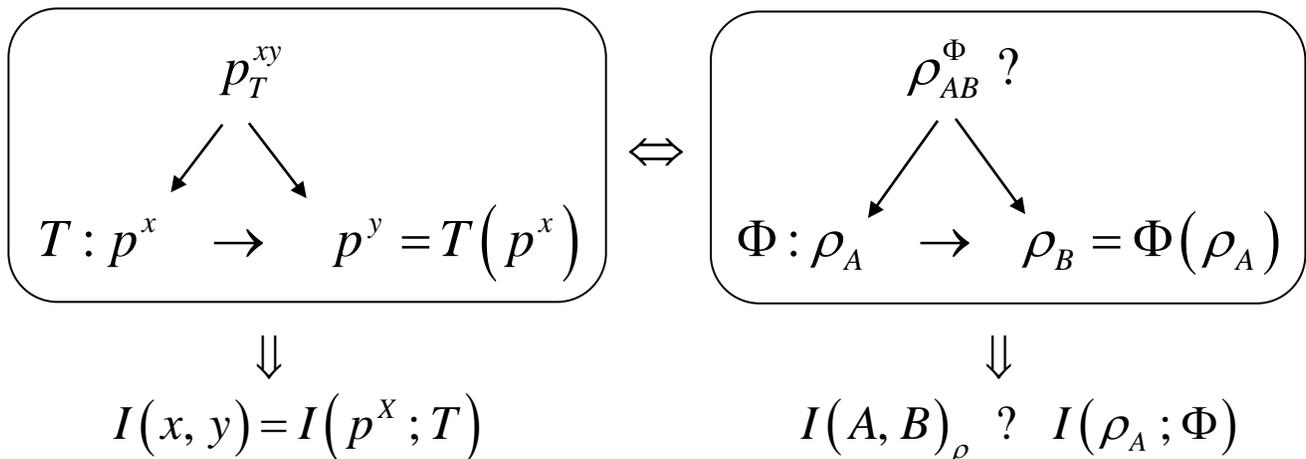
Quantum mutual information via channel Φ in the context of quantum data processing

Entropy exchange (by B. Schumacher):

$$S(\rho_A, \Phi) \equiv S(\rho_E)$$

where $\rho_E = \Phi^c(\rho_A) = \text{Tr}_B V \rho_A V^*$. One can define

$$I(\rho_A; \Phi) \equiv S(\rho_A) + S(\Phi(\rho_A)) - S(\rho_A, \Phi)$$



How can we connect scenario 1 and scenario 2?

Quantum conditional probability and quantum Bayse formula

Quantum conditional probability (by Cref and Adami)

$$\rho_{B|A} = \lim_{n \rightarrow \infty} \left[\rho_{AB}^{1/n} \left(\rho_A^{-1/n} \otimes I_B \right) \right]^n$$

then, we have

$$S(B|A)_\rho \triangleq S(\rho_{AB}) - S(\rho_A) = -\text{Tr} \rho_{AB} \log \rho_{B|A}$$

In this talk:

another two types of quantum conditional probability
(QCP for short)

< **QCP type 1**: full quantum (by Asorey et.al. or Leifer) >

For a given ρ_A and Φ ,

QCP via Choi-Jamiolkowski isomorphism $\pi_{B|A}^{q-q}$:

$$\pi_{B|A}^{q-q} \equiv \sum_{i,j} E_{ij} \otimes \Phi(E_{ij})$$

Compound state ρ_{AB}^{q-q} :

$$\rho_{AB}^{q-q} = \left(\rho_A^{1/2} \otimes I_B \right) \pi_{B|A}^{q-q} \left(\rho_A^{1/2} \otimes I_B \right)$$

where $E_{ij} = |e_i\rangle\langle e_j|$ and $|e_i\rangle$ is the eigenbasis of ρ_A , i.e.,

$$\rho_A = \sum_i p_i |e_i\rangle\langle e_i| = \sum_i p_i \Pi_i$$

then,

$$(1) \quad \pi_{B|A}^{q-q} \geq 0 \quad (2) \quad \text{Tr}_B \pi_{B|A}^{q-q} = I_A$$

Quantum Bayse formula

$$\rho_{AB}^{q-q} = \sum_{i,j} \sqrt{p_i} \sqrt{p_j} E_{ij} \otimes \Phi(E_{ij})$$

$$\Phi: \rho_A = \text{Tr}_B \rho_{AB}^{q-q} \longrightarrow \rho_B = \text{Tr}_A \rho_{AB}^{q-q} = \Phi(\rho_A)$$

Theorem 1

$$I(A, B)_\rho = I(\rho_A; \Phi)$$

Proof :

$$\begin{aligned} I(A, B)_\rho &= S(\rho_A) + S(\rho_B) - S(\rho_{AB}^{q-q}) \\ &= S(\rho_A) + S(\Phi(\rho_A)) - S(\rho_{AB}^{q-q}) \end{aligned}$$

we can show $S(\rho_{AB}^{q-q}) = S(\rho_A : \Phi)$, therefore

$$\begin{aligned} &= S(\rho_A) + S(\Phi(\rho_A)) - S(\rho_A, \Phi) \\ &= I(\rho_A; \Phi) \end{aligned}$$

The quantum version of the classical equality

$$I(x, y) = I(p^x; T)$$

Notice:

In general $I(A, B)_\rho$ or $I(\rho_A; \Phi)$ violates the Shannon's fundamental inequality, for example in the case of pure entangled state.

We rewrite $I(\rho_A, \Phi)$ by $I^{\varrho-\varrho}(\rho_A; \Phi)$

< QCP type 2: semi quantum (by Ohya) >

For a given ρ_A and Φ

Compound state ρ_{AB}^0 :

$$\rho_{AB}^0 \equiv \sum_i p_i E_{ii} \otimes \Phi(E_{ii}) = \sum_i p_i \Pi_i \otimes \Phi(\Pi_i) =: \rho_{AB}^{c-q}$$

$$\rho_A = \sum_i p_i |e_i\rangle\langle e_i| = \sum_i p_i \Pi_i$$

QCP $\pi_{B|A}^{c-q}$: $\pi_{B|A}^{c-q} \equiv \sum_i E_{ii} \otimes \Phi(E_{ii}) = \sum_i \Pi_i \otimes \Phi(\Pi_i)$

↓

$$\rho_{AB}^{c-q} = \left(\rho_A^{1/2} \otimes I_B \right) \pi_{B|A}^{c-q} \left(\rho_A^{1/2} \otimes I_B \right)$$

Proposition 1 The quantum analog of conditional entropy

$$-\text{Tr} \rho_{AB}^{c-q} \log \pi_{B|A}^{c-q} = \sum_i p_i S(\rho_{B|i}) \quad (\Leftrightarrow \sum_i p_i^x S(p_i^y))$$

where $\rho_{B|i} = \Phi(\Pi_i)$ (\leftrightarrow conditional state).

The mutual entropy via Φ : scenario 3 (by Ohya)

$$I^0(p_A; \Phi) \equiv \sup_{\Pi_i} D(\rho_{AB}^{c-q} \| \rho_A \otimes \Phi(\rho_A)) =: I^{c-q}(\rho_A; \Phi)$$

where the sup is taken all decomposition $\rho_A = \sum_i p_i \Pi_i$.

Proposition 2

$$I^{c-q}(\rho_A; \Phi) = S(\Phi(\rho_A)) - \inf_{\Pi_i} \sum_k p_k S(\Phi(\Pi_k))$$

$$= \sup_{\Pi_i} \sum_k p_k D(\Phi(\Pi_k) \| \Phi(\rho_A))$$

Theorem 2 $0 \leq I^{c-q}(p_A; \Phi) \leq \min \{ S(\rho_A), S(\Phi(\rho_A)) \}$

< The relation between $I(A, B)_\rho = I^{Q-Q}(\rho_A; \Phi)$
and $I^{C-Q}(\rho_A; \Phi)$ >

Proposition 3

Let Π be the quantum channel defined by

$$\Pi(\rho) \equiv \sum_i \Pi_i \rho \Pi_i.$$

One has

$$I^{C-Q}(\rho_A; \Phi) = \sup_{\Pi_i} I^{Q-Q}(\rho_A; \Phi_{\Pi})$$

where $\Phi_{\Pi} = \Phi \circ \Pi$, and the sup is taken all decomposition
 $\rho_A = \sum_i p_i \Pi_i$.

Proposition 4 One has the following relation

$$D(\rho_{AB}^{q-q} \parallel \rho_A \otimes \rho_B) = D(\rho_{AB}^{q-q} \parallel \rho_{AB}^{c-q}) + D(\rho_{AB}^{c-q} \parallel \rho_A \otimes \rho_B)$$

and hence

$$I^{Q-Q}(\rho_A; \Phi) = \inf_{\Pi_i} D(\rho_{AB}^{q-q} \parallel \rho_{AB}^{c-q}) + I^{C-Q}(\rho_A; \Phi)$$

Note: the difference

$$Q(\rho_A; \Phi) := I^{Q-Q}(\rho_A; \Phi) - I^{C-Q}(\rho_A; \Phi) = \inf_{\Pi_i} D(\rho_{AB}^{q-q} \parallel \rho_{AB}^{c-q})$$

defines relative entropy distance between ρ_{AB}^{q-q} and the
closest ρ_{AB}^{c-q} with the same marginal states ρ_A and
 $\rho_B = \Phi(\rho_A) = \Phi_{\Pi}(\rho_A)$.

The quantum discord vs $Q(\rho_A; \Phi)$

The notion of quantum discord:

(by Ollivier & Zurek or Henderson & Vedral)

Step1 : a measurement on A part by the correction of one-dimensional projectors Π_k in \mathbb{H}_A satisfying $\sum_k \Pi_k = I_A$,

$$\rho_{AB} \rightarrow \rho_{AB}^{C-Q} := [\Pi \otimes \text{id}_B] \rho_{AB} = \sum_k (\Pi_k \otimes I_B) \rho_{AB} (\Pi_k \otimes I_B)$$

Step2 : the state change after the measurement with outcome k

$$k \rightarrow \rho_{AB|k} \equiv \frac{1}{p_k} (\Pi_k \otimes I_B) \rho_{AB} (\Pi_k \otimes I_B) \rightarrow \rho_{B|k} = \text{Tr}_A \rho_{AB|k}$$

where $p_k = \text{Tr}(\Pi_k \otimes I_B) \rho_{AB} (\Pi_k \otimes I_B) = \text{Tr}(\Pi_k \rho_A)$.

< Quantum discord - scenario 4 ->

$$D(\rho_{AB}) \equiv I(A:B)_\rho - J(B|A)$$

where $J(B|A) \equiv \sup_{\{\Pi_k\}} I(\rho_{AB} | \{\Pi_k\})$ and

$$I(\rho_{AB} | \{\Pi_k\}) \equiv S(\rho_B) - \sum_k p_k S(\rho_{B|k}) = D(\rho_{AB}^{C-Q} \| \rho_A \otimes \rho_B)$$

where

$$\rho_{AB}^{C-Q} = \sum_k p_k \Pi_k \otimes \rho_{B|k}$$

with $\rho_A = \sum_k p_k \Pi_k$ and $\rho_B = \sum_k p_k \rho_{B|k}$

$D(\rho_{AB})$ may be equivalently $Q(\rho_A; \Phi)$

Let replace ρ_{AB} by ρ_{AB}^{q-q} , then

$$\begin{aligned} D(\rho_{AB}^{q-q}) &= \inf_{\Pi_i} \left\{ I(A, B)_{\rho} - I(A, B)_{\rho^{c-q}} \right\} \\ &= \inf_{\Pi_i} \left\{ I^{\mathcal{Q}-\mathcal{Q}}(\rho_A; \Phi) - D(\rho_{AB}^{c-q} \| \rho_A \otimes \Phi(\rho_A)) \right\} \\ &= I^{\mathcal{Q}-\mathcal{Q}}(\rho_A; \Phi) - I^{c-\mathcal{Q}}(\rho_A; \Phi) \\ &= Q(\rho_A; \Phi) \end{aligned}$$

Notice:

In the case of ρ_{AB}^{q-q} the rank-1 projectors have to commute with ρ_A .

Such quantity may be called the “**measurement-induced disturbance**” (by Luo)

= One performs a projective measurement on the A subsystem without perturbing its system. =

Definition 1 The measurement-induced disturbance of the channel Φ is defined by

$$\Delta(\Phi) \equiv \inf_{\rho_A} Q(\rho_A; \Phi)$$

$\Delta(\Phi)$ measures the distortion of Φ from the c-q channel Φ^{c-q} due to the fact:

$$\Delta(\Phi^{c-q}) = 0$$

Summary

Introducing the quantum analog of Bayse formula via QCP (full quantum and semi quantum) , we made clear the relation between several types of mutual entropy.

Type 1 (full quantum QCP $\pi_{B|A}^{q-q}$)

$$\rho_{AB}^{q-q} = (\rho_A^{1/2} \otimes I_B) \left(\sum_i E_{ij} \otimes \Phi(E_{ij}) \right) (\rho_A^{1/2} \otimes I_B)$$

we have

$$\begin{aligned} I(A, B)_\rho &= S(\rho_A) + S(\rho_B) - S(\rho_{AB}^{q-q}) \\ &= S(\rho_A) + S(\Phi(\rho_A)) - S(\rho_A, \Phi) = I^{\varrho-\varrho}(\rho_A; \Phi) \end{aligned}$$

Type 2 (semi quantum QCP $\pi_{B|A}^{c-q}$)

$$\rho_{AB}^{c-q} = (\rho_A^{1/2} \otimes I_B) \left(\sum_i \Pi_i \otimes \Phi(\Pi_i) \right) (\rho_A^{1/2} \otimes I_B)$$

we have

$$\begin{aligned} I^{c-\varrho}(\rho_A; \Phi) &= S(\Phi(\rho_A)) - \inf_{\Pi_i} \sum_k p_k S(\Phi(\Pi_k)) \\ 0 \leq I^{c-\varrho}(\rho_A; \Phi) &\leq \min \{ S(\rho_A), S(\Phi(\rho_A)) \} \end{aligned}$$

$$I^{c-\varrho}(\rho_A; \Phi) = \sup_{\Pi_i} I^{\varrho-\varrho}(\rho_A; \Phi_{\Pi}) \quad (\Phi_{\Pi} = \Phi \circ \Pi)$$

Two types of QCP define

the measurement-induced disturbance of the channel Φ

$$Q(\rho_A; \Phi) := I^{\varrho-\varrho}(\rho_A; \Phi) - I^{c-\varrho}(\rho_A; \Phi) = \inf_{\Pi_i} D(\rho_{AB}^{q-q} \parallel \rho_{AB}^{c-q})$$

$$\Delta(\Phi) = \inf_{\rho_A} Q(\rho_A; \Phi)$$

Thank you for your attention!!